



Monte Carlo Methods

Spring Semester 2013/14, Department of Applied Mathematics, University of Crete

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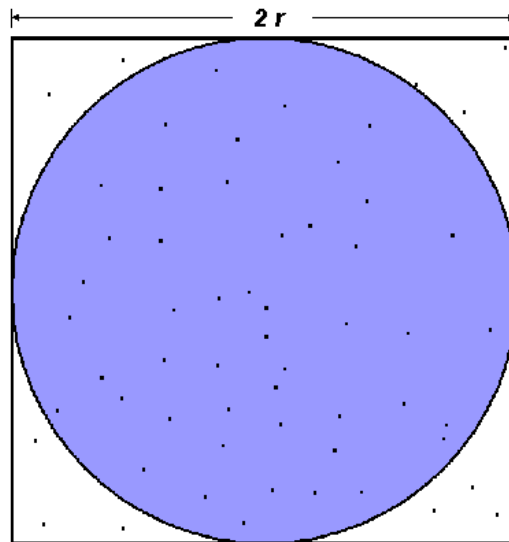
**Part II: Monte Carlo Integration, Random Numbers
generators**



Introductory Examples: Calculate π

Calculation of number π with the following method:

- Περικλείουμε κύκλο με ένα τετράγωνο. Δημιουργούμε m τυχαία σημεία μέσα στο τετράγωνο.
- Βρίσκουμε τα σημεία που εμπεριέχονται και μέσα στον κύκλο, n .
- Αν $r = n/m$, τότε ο αριθμός π προσεγγίζεται ως $\pi \approx 4r$. Όσο περισσότερα τα σημεία m τόσο μεγαλύτερη ακρίβεια του υπολογισμού.



$$A_S = (2r)^2 = 4r^2$$

$$A_C = \pi r^2$$

$$\pi = 4 \times \frac{A_C}{A_S}$$



Introductory Examples: Calculate π

Algorithm:

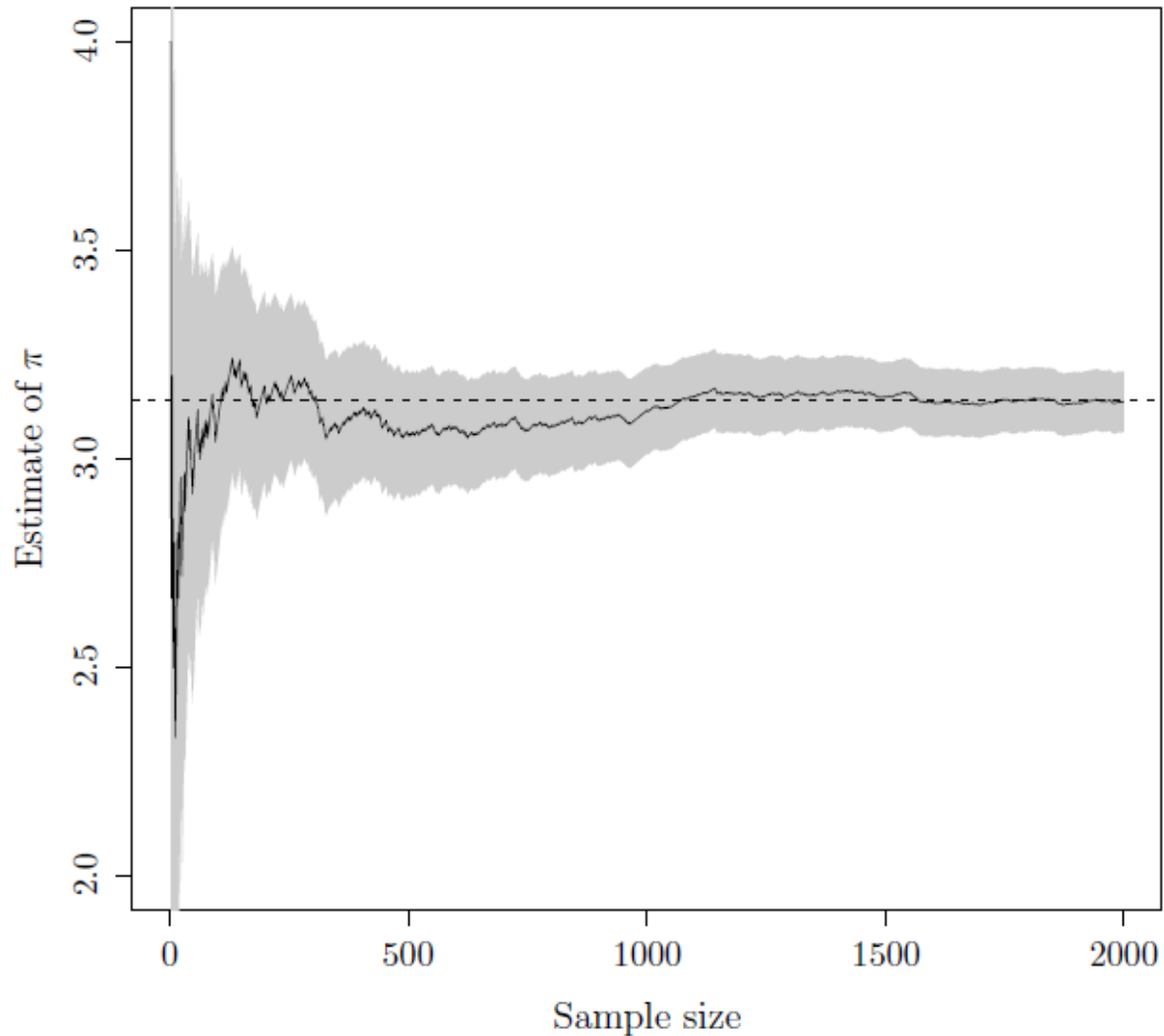
```
npoints = 1000000
circle_count = 0
do j = 1, npoints
    generate 2 random numbers between 0 and 1
    xcoordinate = random1
    ycoordinate = random2
    if (xcoordinate, ycoordinate) inside circle then
        circle_count = circle_count + 1
    end do
end do
PI = 4.0*circle_count/npoints
```

- Ο χρόνος υπολογισμού είναι κυρίως ο χρόνος εκτέλεσης της επαναληπτικής διαδικασίας (loop).
- Αυτό οδηγεί σε (σχεδόν) ‘τέλειο παραλληλισμό’ (**embarrassingly parallelism**):
 - Εντατικοί υπολογισμοί.
 - Ελάχιστη επικοινωνία, ελάχιστο I/O.



Introductory Examples: Calculate π

□ Estimate π as a function of sample size:





Monte Carlo Integration

- Two major classes of numerical problems that arise in statistical inference
 - *optimization* problems
 - *integration* problems
- Although optimization is generally associated with the likelihood approach, and integration with the Bayesian approach, these are not strict classifications

- Generic problem of evaluating the integral

$$E_f[h(X)] = \int_{\mathcal{X}} h(x) f(x) dx .$$

- Based on previous developments, it is natural to propose using a sample (X_1, \dots, X_m) generated from the density f
- Approximate the integral by the empirical average
- This approach is often referred to as the *Monte Carlo method*



Monte Carlo Integration

Strong Law

- For a sample (X_1, \dots, X_m) , the empirical average

$$\bar{h}_m = \frac{1}{m} \sum_{j=1}^m h(x_j) ,$$

converges almost surely to

$$E_f[h(X)]$$

- This is the Strong Law of Large Numbers



Monte Carlo Integration

Central Limit Theorem

- Estimate the variance with

$$\text{var}(\bar{h}_m) = \frac{1}{m} \int_{\mathcal{X}} (h(x) - E_f[h(X)])^2 f(x) dx$$

- For m large,

$$\frac{\bar{h}_m - E_f[h(X)]}{\sqrt{v_m}}$$

is therefore approximately distributed as a $\mathcal{N}(0, 1)$ variable

- This leads to the construction of a convergence test and of confidence bounds on the approximation of $E_f[h(X)]$.



Monte Carlo Integration: Example

□ Example: Calculate the integral of a function $h(x)$

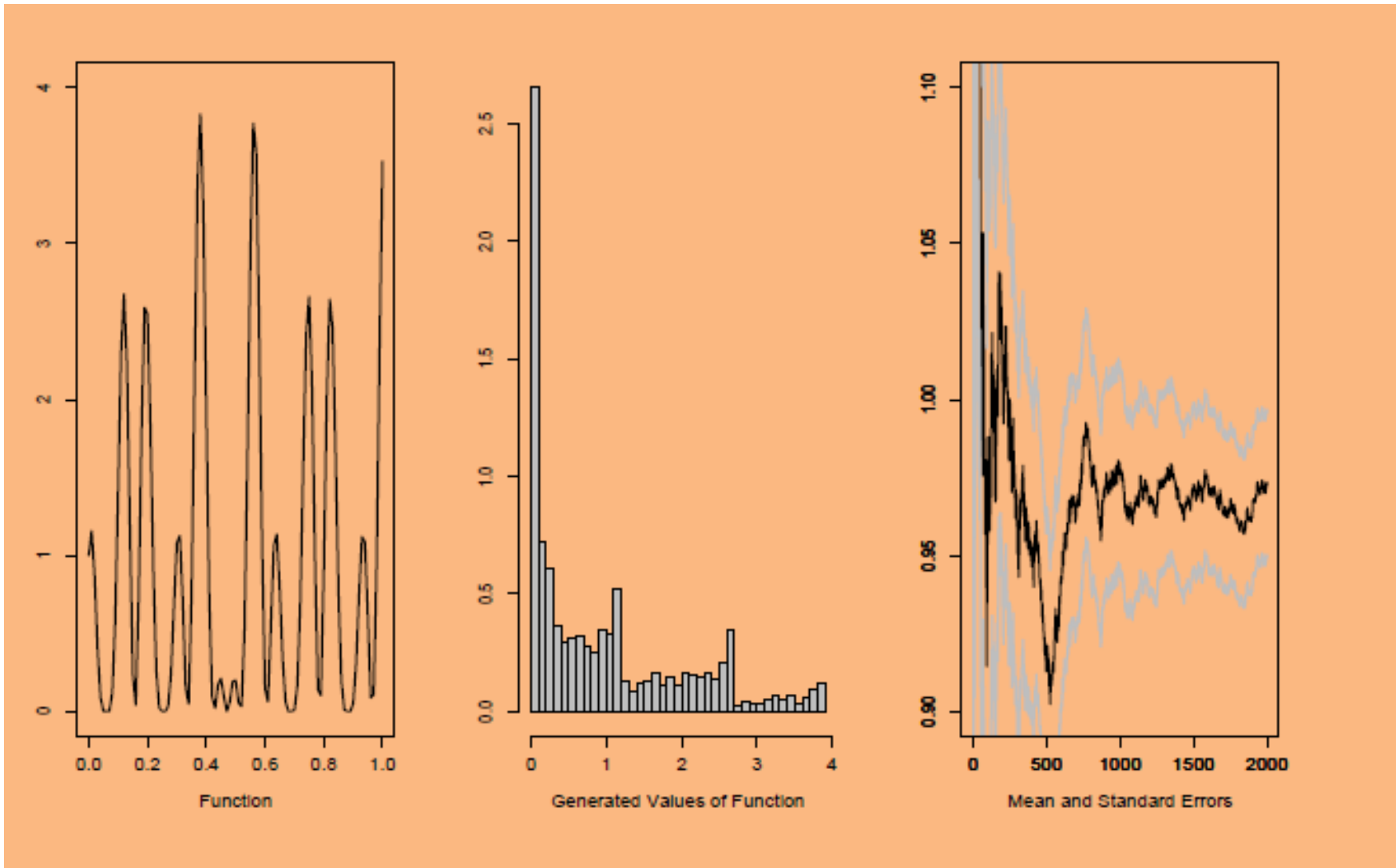
$$h(x) = [\cos(50x) + \sin(20x)]^2$$

- To calculate the integral, we generate U_1, U_2, \dots, U_n iid $\mathcal{U}(0, 1)$ random variables, and approximate $\int h(x)dx$ with $\sum h(U_i)/n$.
- It is clear that the Monte Carlo average is converging, with value of 0.963 after 10,000 iterations.



Monte Carlo Integration: Example

□ Example: Estimators

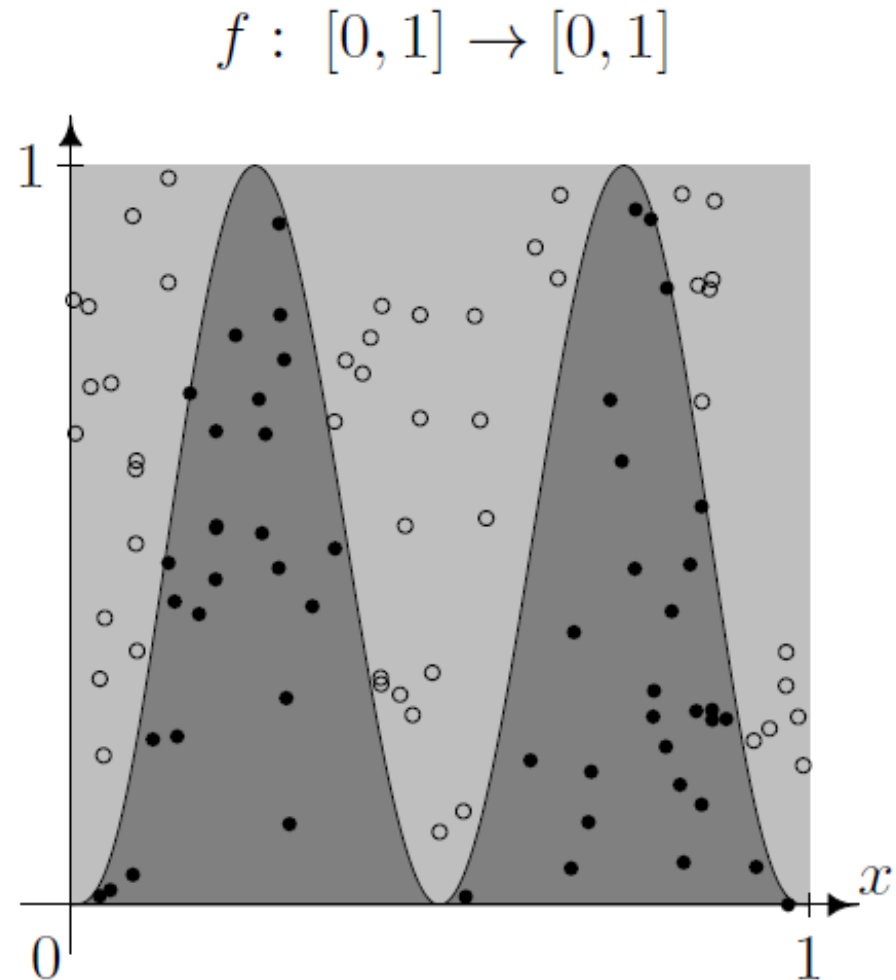




Monte Carlo Integration

- Generalization of Integration: **Riemann sums vs MC method** (see hand notes).

$$\begin{aligned} & \int_0^1 f(x) dx \\ = & \int_0^1 \int_0^{f(x)} 1 dt dx \\ = & \iint_{\{(x,t):t \leq f(x)\}} 1 dt dx \\ = & \frac{\iint_{\{(x,t):t \leq f(x)\}} 1 dt dx}{\iint_{\{0 \leq x,t \leq 1\}} 1 dt dx} \end{aligned}$$





Monte Carlo Integration

□ Comparison – Speed of Convergence:

- Speed of convergence of Monte Carlo integration is $O_{\mathbb{P}}(n^{-1/2})$.
- Speed of convergence of numerical integration of a *one-dimensional* function by Riemann sums is $O(n^{-1})$.
- Does not compare favourably for one-dimensional problems.
- However:
 - Order of convergence of Monte Carlo integration is *independent* of the dimension.
 - Order of convergence of numerical integration techniques like Riemann sums deteriorates with the dimension increasing.

↪ Monte Carlo methods can be a good choice for high-dimensional integrals.



Random Number Generators

- Philosophical paradox:
 - We need to reproduce randomness by a computer algorithm.
 - A computer algorithm is deterministic in nature.

↪ “pseudo-random numbers”
- Pseudo-random number from $U[0, 1]$ will be our only “source of randomness” .
- Other distributions can be derived from $U[0, 1]$ pseudo-random numbers using deterministic algorithms.





Pseudo-Random Number Generators

- A pseudo-random number generator (RNG) should produce output for which the $U[0, 1]$ distribution is a suitable model.
- The pseudo-random numbers X_1, X_2, \dots should thus have the same *relevant* statistical properties as independent realisations of a $U[0, 1]$ random variable.
 - They should reproduce independence (“lack of predictability”): X_1, \dots, X_n should not contain any discernible information on the next value X_{n+1} . This property is often referred to as the lack of predictability.
 - The numbers generated should be spread out evenly across $[0, 1]$.



Pseudo-Random Number Generators

□ A simple example: Congruential pseudo-RNG.

Algorithm 1.1: Congruential pseudo-random number generator

1. Choose $a, M \in \mathbb{N}$, $c \in \mathbb{N}_0$, and the initial value (“seed”) $Z_0 \in \{1, \dots, M - 1\}$.
2. For $i = 1, 2, \dots$
Set $Z_i = (aZ_{i-1} + c) \bmod M$, and $X_i = Z_i/M$.

$Z_i \in \{0, 1, \dots, M - 1\}$, thus $X_i \in [0, 1)$.



Pseudo-Random Number Generators

Consider the choice of $a = 81$, $c = 35$, $M = 256$, and seed $Z_0 = 4$.

$$Z_1 = (81 \cdot 4 + 35) \bmod 256 = 359 \bmod 256 = 103$$

$$Z_2 = (81 \cdot 103 + 35) \bmod 256 = 8378 \bmod 256 = 186$$

$$Z_3 = (81 \cdot 186 + 35) \bmod 256 = 15101 \bmod 256 = 253$$

...

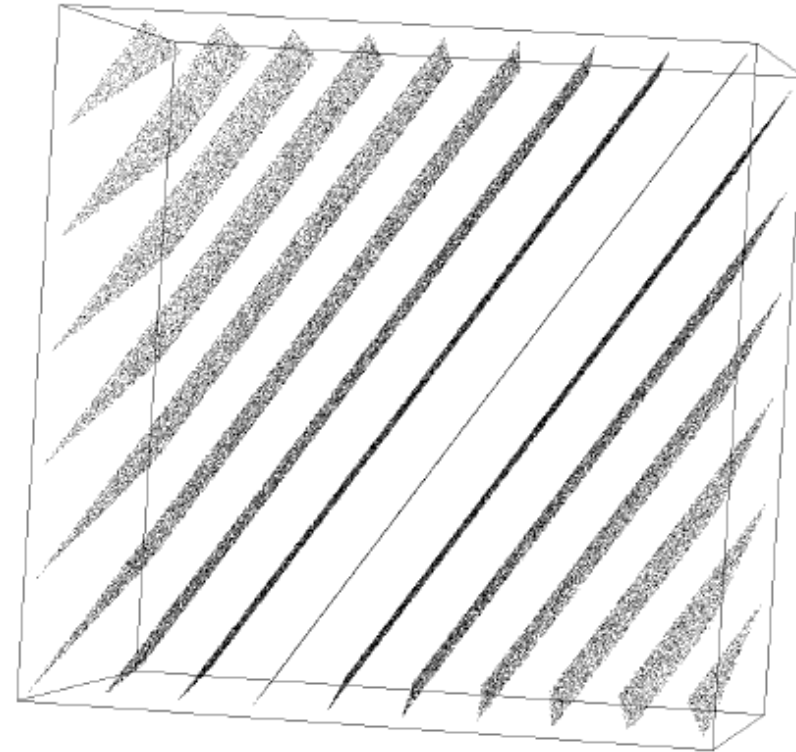
The corresponding X_i are $X_1 = 103/256 = 0.4023438$,
 $X_2 = 186/256 = 0.72656250$, $X_3 = 253/256 = 0.98828120$.



Pseudo-Random Number Generators

❑ RANDU: A typical poor choice of RNG.

- Very popular in the 1970s (e.g. System/360, PDP-11).
- Linear congruential generator with $a = 2^{16} + 3$, $c = 0$, and $M = 2^{31}$.
- The numbers generated by RANDU lie on only 15 hyperplanes in the 3-dimensional unit cube!



According to a salesperson at the time: “We guarantee that each number is random individually, but we don’t guarantee that more than one of them is random.”



Pseudo-Random Number Generators

❑ Flaw of the linear congruential RNG.

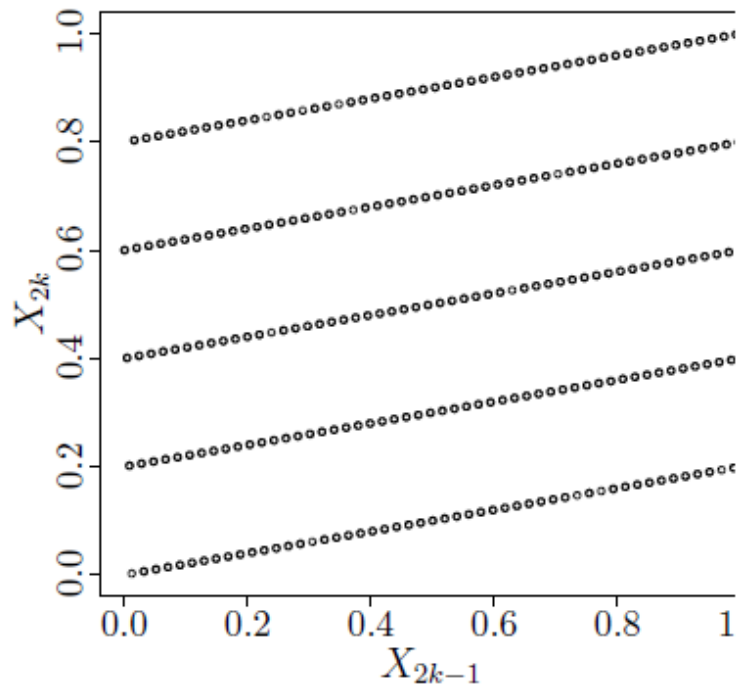
- “Crystalline” nature is a problem for every linear congruential generator.
- Sequence of generated values X_1, X_2, \dots viewed as points in an n -dimension cube lies on a finite, and often very small number of parallel hyperplanes.
- Marsaglia (1968): “the points [generated by a congruential generator] are about as randomly spaced in the unit n -cube as the atoms in a perfect crystal at absolute zero.”
- The number of hyperplanes depends on the choice of a , c , and M .
- For these reasons **do not use the linear congruential generator!** Use more powerful generators (like e.g. the *Mersenne twister*, available in GNU R).



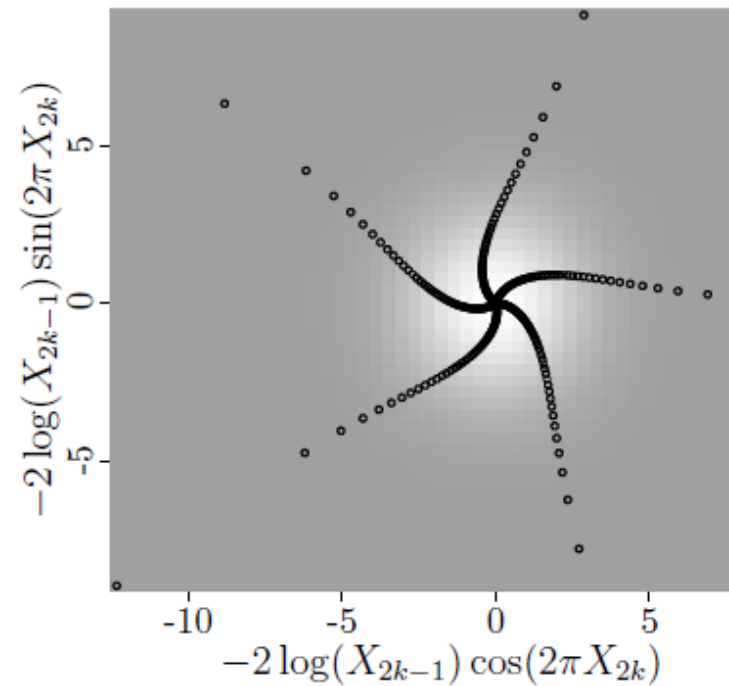
Pseudo-Random Number Generators

□ Another problematic example:

Linear congruential generator with $a = 1229$, $c = 1$, and $M = 2^{11}$.



Pairs of generated values (X_{2k-1}, X_{2k})



Transformed by Box-Muller method



Bibliography

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